

A Compositional Semantics for Venn Diagrams

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The important work of [3], [2], and others has applied techniques for analyzing formal languages to Venn Diagrams. This work has shown that Venn Diagrams can be specified by recursive formation rules, given a model-theoretic semantics, and a complete proof theory. But these theories are not rule-to-rule compositional. We will develop a rule-to-rule compositional semantics for Venn Diagrams. This endeavor sheds light on whether Venn diagrams present a distinct mode of representation compared to language, and the extent to which “the elements and combinatorial rules for [Venn] diagrams are very different than those for sentences” ([1]: 153-4).

Our investigation delineates two methods for syntactically deriving a Venn Diagram. The *cell-first* approach first constructs the cells or minimal regions of a diagram out of its basic regions. Notice that diagram 1 has eight cells, *a-h*.

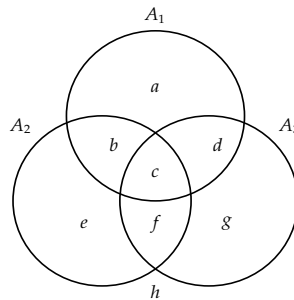


Diagram 1: Undecorated Venn diagram with labeled cells

The cell-first approach then adds markers to indicate that the region represents a non-empty or empty set. The *region-first approach*, by contrast, directly constructs non-basic regions and adds markers to indicate that the region represents an empty or non-empty set. Thus, to indicate that the shaded region represents an empty set, the former would construct Diagram 2 in two steps by shading cell *b* and then cell *c*, while the latter would directly shade the region $b \cup c$ to indicate that it represents an empty set. Likewise for “salvation”, such as in Diagram 3, where hatch marks indicate that the region represents a non-empty set.

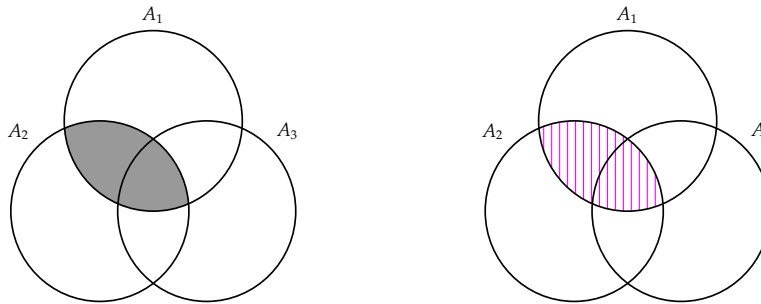


Diagram 2: A diagram with region $b \cup c$ destroyed
 Diagram 3: A diagram with region $b \cup d$ saved

In this talk, we focus on the region-first approach. Let a diagram $D = \langle B, \Delta, \Sigma \rangle$ where B is the set of basic regions; Δ is the set of destroyed regions; and Σ is the set of saved regions. Venn Diagrams can be constructed via the following syntactic construction rules:

Basic Regions: $\{A_1, A_2, \dots, A_i, \dots\}$

Form Rule: If X_1, \dots, X_n are basic regions satisfying the mutual overlap constraint, then $\text{form}(X_1, \dots, X_n)$ is a diagram with regions including X_1, \dots, X_n and R .

Intersection Rule: If X_1 and X_2 are regions of D , then $\text{inter}(X_1, X_2) = \bigcap\{X_1, X_2\}$ is a region of D , provided that it is non-empty.

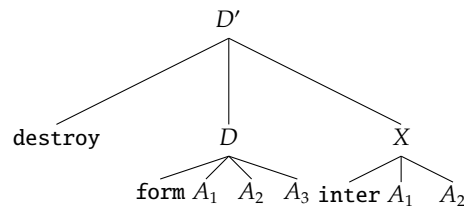
Union Rule: If X_1 and X_2 are regions of D , then $\text{union}(X_1, X_2) = \bigcup\{X_1, X_2\}$ is a region of D , provided that it is non-empty.

Complement Rule: If X_1 and X_2 are regions of D , then $\text{comp}(X_1, X_2) = X_1 - X_2$ is a region of D , provided that it is non-empty.

Destruction Rule: If D is a diagram and if X is a region of D , then $\text{destroy}(D, X)$ is a diagram.

Salvation Rule: If D is a diagram and if X is a region of D , then $\text{save}(D, X)$ is a diagram.

For example this is the syntactic derivation of Diagram 2:



That's the syntax.³ Now the clauses for the compositional semantics.

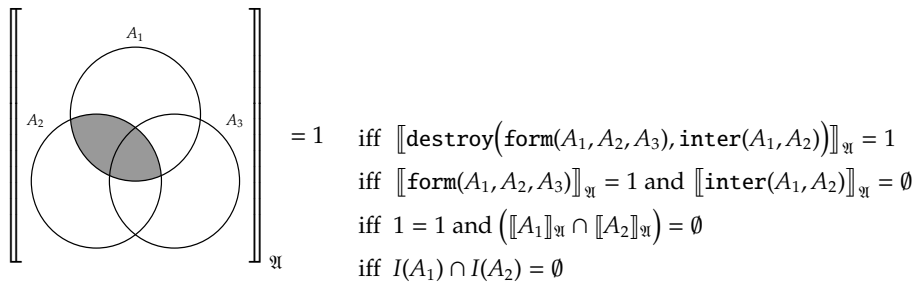
– Semantics for Regions

- For any basic region X , $\llbracket X \rrbracket_{\mathfrak{A}} = I(X)$, where $\mathfrak{A} = \langle U, I \rangle$.
- $\llbracket \text{inter}(X_1, X_2) \rrbracket_{\mathfrak{A}} = \llbracket X_1 \rrbracket_{\mathfrak{A}} \cap \llbracket X_2 \rrbracket_{\mathfrak{A}}$
- $\llbracket \text{union}(X_1, X_2) \rrbracket_{\mathfrak{A}} = \llbracket X_1 \rrbracket_{\mathfrak{A}} \cup \llbracket X_2 \rrbracket_{\mathfrak{A}}$
- $\llbracket \text{comp}(X_1, X_2) \rrbracket_{\mathfrak{A}} = \llbracket X_1 \rrbracket_{\mathfrak{A}} - \llbracket X_2 \rrbracket_{\mathfrak{A}}$

– Semantics for Diagrams

- $\llbracket \text{form}(X_1, \dots, X_n) \rrbracket_{\mathfrak{A}} = 1$.
- $\llbracket \text{destroy}(D, X) \rrbracket_{\mathfrak{A}} = \begin{cases} 1, & \text{if } \llbracket D \rrbracket_{\mathfrak{A}} = 1 \text{ and } \llbracket X \rrbracket_{\mathfrak{A}} = \emptyset, \\ 0, & \text{otherwise.} \end{cases}$
- $\llbracket \text{save}(D, X) \rrbracket_{\mathfrak{A}} = \begin{cases} 1, & \text{if } \llbracket D \rrbracket_{\mathfrak{A}} = 1 \text{ and } \llbracket X \rrbracket_{\mathfrak{A}} \neq \emptyset, \\ 0, & \text{otherwise.} \end{cases}$

The truth-conditions of a diagram such as Diagram 2 = $\text{destroy}(\text{form}(A_1, A_2, A_3), \text{inter}(A_1, A_2))$ can then be derived as follows.



References

1. Camp, E.: Thinking with maps. *Philosophical Perspectives* **21**(1), 145–182 (2007). <https://doi.org/10.1111/j.1520-8583.2007.00124.x>
2. Hammer, E., Danner, N.: Towards a model theory of diagrams. *Journal of Philosophical Logic* **25**(5), 463–482 (1996). <https://doi.org/10.1007/BF00257381>
3. Shin, S.J.: *The logical status of diagrams*. Cambridge University Press (1994)

³ Provided a definition of the *mutual overlap constraint*. A (non-empty) set of basic regions $\{X_1, \dots, X_n\}$ of the plane R meets the mutual overlap constraint iff for all Γ : if $\Gamma \subseteq \{X_1, \dots, X_n\}$, then $\bigcap (\Gamma \cup \{R\}) - \bigcup (\{X_1, \dots, X_n\} - \Gamma) \neq \emptyset$.